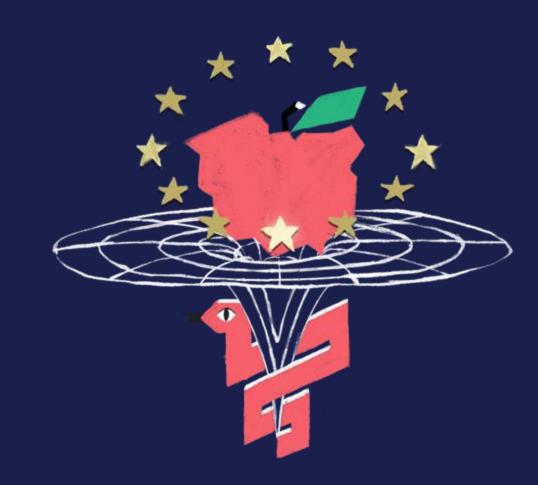


Solving Hard Optimization Problems with Pyomo and HiGHS

A Practical Introduction



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- Modern Data Warehousing & Analytics
- Personalisation & RecSys
- Uncertainty Quantification & Causality
- Python Data Stack
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Definition of Linear Programming (LP)

minimize
$$m{c}^T m{x}$$
 subject to $m{A} m{x} \leq m{b}$ $m{x} \in \mathbb{R}^n$

LPs can be solved with the Simplex algorithm in cubic on average but exponential number of steps in worst case scenario or interior point methods in about $O(n^{2.5})$.



Definition of Mixed Integer Programming (MIP)

$$egin{aligned} \mathsf{minimize} & oldsymbol{c}^Toldsymbol{x} \ \mathsf{subject\ to} & oldsymbol{A}oldsymbol{x} \leq oldsymbol{b} \ oldsymbol{x} \in \mathbb{R}^n \ & x_j \in \mathbb{Z}\ \mathsf{for}\ j \in J \end{aligned}$$

MIPs are NP-hard, i.e. complexity grows exponentially with *n* Dropping the integrality constraints, i.e. "linear relaxation", leads to an LP problem.



Special Cases of Mixed Integer Programming

The cost functional f(x) can be

- linear, i.e. $c^T x$,
- quadratic, i.e. $x^TQx + q^Tx$,

• . . .

If all variables need to be integer, it is a (pure) integer linear program (ILP, IP).

If all variables need to be 0 or 1 (binary, boolean), it is a **0 - 1 linear program**.





Typical fields of application

- Assignment/Allocation Problem

 How to schedule some talks in an agenda under some constraints like room sizes to optimize the conference experience?
- Transportation Problem
 Which supplier should deliver to which factories given some costs to satisfy each factory with minimal costs?
- Shortest Path Problem
 Given some graph with a cost on each edge, what is the shortest path from source to sink?
- Maximum Flow Problem
 Given some pipe system with a source and sink, what is the maximum flow through the system?



Solving a MILP

Two major method classes for solving MILPs:

- Cutting plane methods
- Branch and bound methods

Often combined as Branch and Cut methods

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Source: Introduction to Optimization by Laurent Lessard, Spring 2017–18



Cutting Planes Methods

Cutting Planes

Idea

- solve the LP relaxation problem
- while solution is not integral:
 - o add constraint that excludes solution but no integer points
 - o solve LP relaxation again

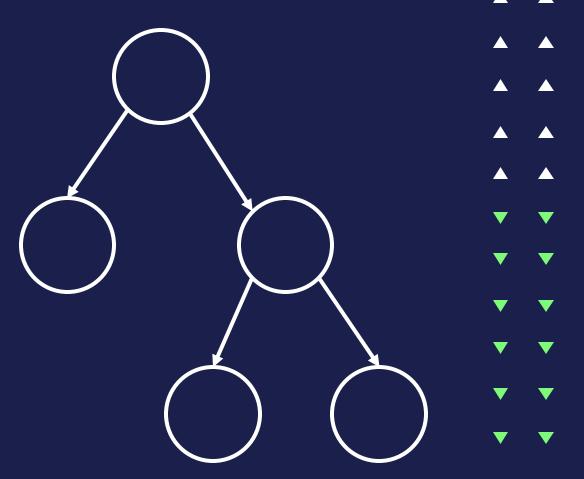


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Branch & Bound Methods

Branch & Bound

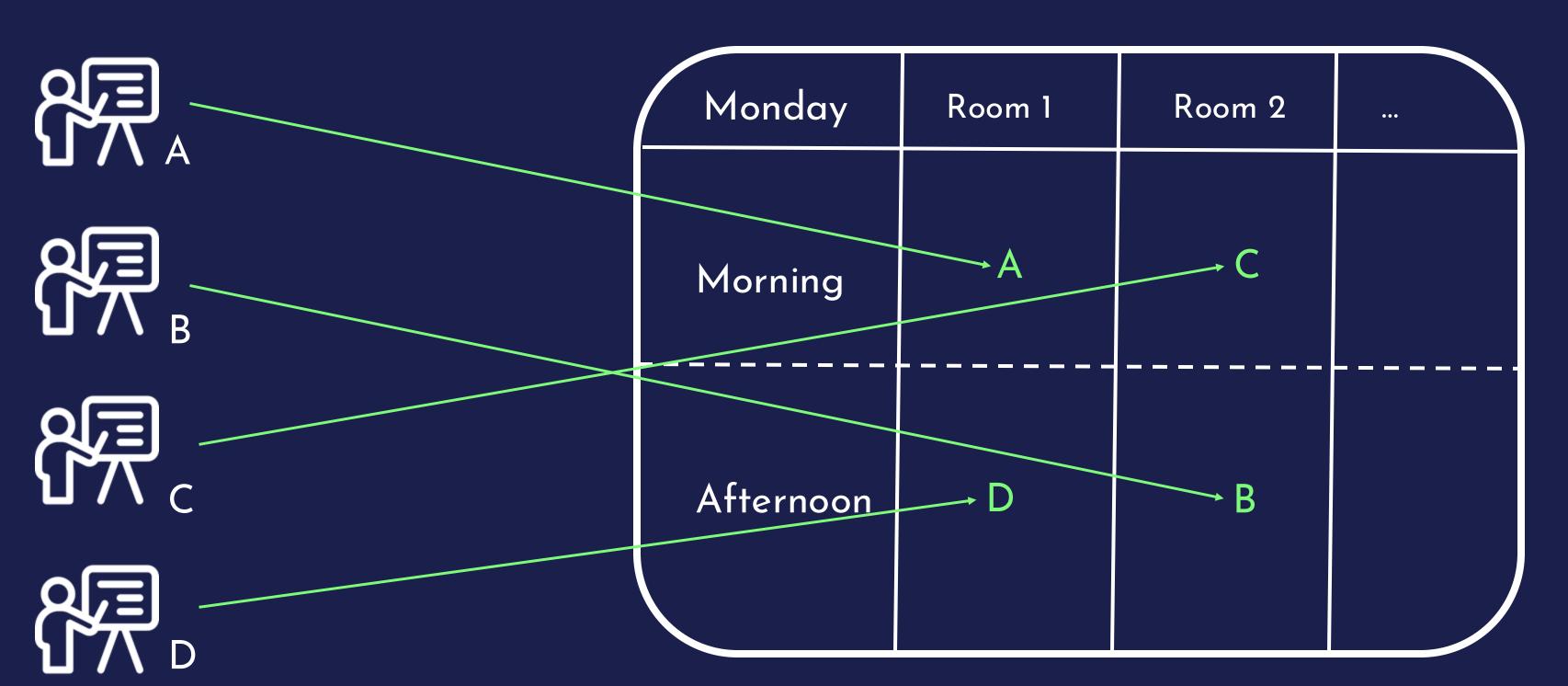


Idea

- 1. solve the relaxed LP and for a fractional x_i split into two subproblems
 - \circ with constraint $x_i \leq \lfloor x_i \rfloor$
 - \circ and with constraint $x_i \geq \lceil x_i \rceil$
- 2. repeat step 1. and build a tree of subproblems
- 3. eliminate branches of the tree using $f(x^{feasible}) \le f(x^{optimal}) \le f(x^{relaxed})$



How can we use MILPs for a Conference Scheduling?





What Constraints do we have?

- each talk must be assigned exactly once,
- each room/timeslot combination can only be occupied by one talk at most,
- the length of the timeslot must match the length of the talk
- some tutorials have part 1 & 2, thus need to be consecutive





What is our objective?

- the preferences for day and time of some speakers, e.g. keynotes, need to be considered
- 2.popularity of a talk should be reflected in the room capacity,
- 3.avoid parallel talks that attract the same audience,
- 4.have in the same session, i.e. consecutive block of talks, the same main track, e.g. PyData vs. PyCon,
- 5.or even the same sub track, e.g. PyData: Data Handling,

Precedence: 1 > 2 > 3 > 4 > 5



Solving MILPs in Python

1. Framework to formulate the problem, e.g. a.Pyomo by Sandia NL



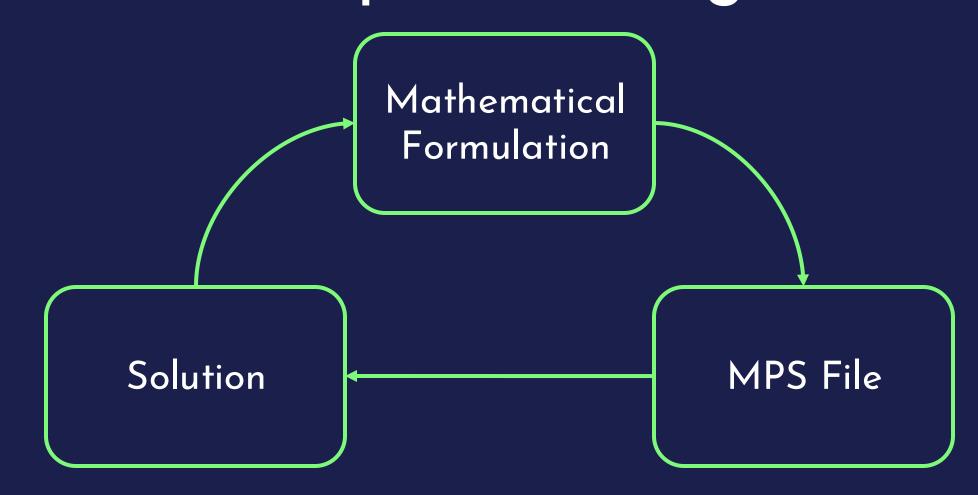
1. Solver to solve the canonical problem, e.g.



b.CBC

b.PulP

c.IPOPT







Pyomo: Index Sets

```
model = pyo.ConcreteModel(name="PyConDE/PyData Schedule")
model.sTalks = pyo.Set(initialize=talks_df[Col.submission].values)
model.sDays = pyo.Set(initialize=["Wednesday", "Thursday", "Friday"])
model.sSessions = pyo.Set(initialize=["Morning", "Afternoon1", "Afternoon2"])
model.sSlots = pyo.Set(initialize=["First", "Second", "Third"])
model.sRooms = pyo.Set(initialize=room_caps_dict.keys())
model.sMainTracks = pyo.Set(initialize=main_tracks)
model.sSubTracks = pyo.Set(initialize=sub_tracks)
```



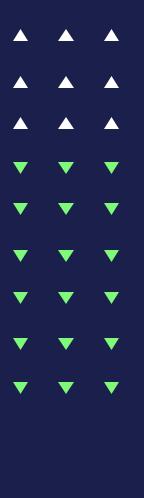
Pyomo: Parameters

```
def init_slot_preference(model, t, d, s, l, r):
    if t in sponsored_talks and r not in {Rooms.PALLADIUM, Rooms.HASSIUM}
      and not (d == "Friday" and s in {"Afternoon1", "Afternoon2"}):
        return 1
    elif t in {"S8MUBF"} and d == "Wednesday" and s == "Morning":
        return 1
    elif t in {"KCV9RS"} and d == "Thursday" and s == "Morning":
        return 1
    elif d == "Friday" and s == "Afternoon2" and l == "Third":
        return -1
    else:
        return 0
model.pPreferences = pyo.Param(model.sTalks, model.sDays, model.sSessions,
                               model.sSlots, model.sRooms, initialize=init_slot_preference)
# Same for pTalkRoomFit, pCoOccurencesPenalty, pSlotLengths, pTalkLengths, ...
```



Pyomo: Variables

```
## Decision variable
model.vbSchedule = pyo.Var(model.sTalks, model.sDays, model.sSessions,
                           model.sSlots, model.sRooms, domain=pyo.Binary)
## Auxiliary variables
# indicator if talk t is in room r
model.vbTalkRoom = pyo.Expression(model.sTalks, model.sRooms,
    rule=lambda model, t, r: sum(model.vbSchedule[t, d, s, l, r]
      for d, s, l in product(model.sDays, model.sSessions, model.sSlots))
# indicator if two talks are in the same day/session/slot combination
model.vbCoOccurences = pyo.Var(model.sTalks, model.sTalks, domain=pyo.Binary)
```





Pyomo: Constraints

```
# Make sure each talk is assigned once
model.ctTalkAssigned = pyo.ConstraintList()
for t in model.sTalks:
    model.ctTalkAssigned.add(sum(model.vbSchedule[t, :, :, :, :]) == 1)

# Make sure each room/timeslot-combination is occupied only with one talk at most
model.ctTimeRoomOccup = pyo.ConstraintList()
for d, s, l, r in product(model.sDays, model.sSessions, model.sSlots, model.sRooms):
    model.ctTimeRoomOccup.add(sum(model.vbSchedule[:, d, s, l, r]) <= 1)

# Also for parameters, e.g. pTalkLengths, and auxilliary variables, e.g. vbCoOccurences.</pre>
```



Pyomo: Objective



Check out the Full Source Code

Pytanis includes a Pretalx client and all the tooling you need for conferences using Pretalx, from handling the initial call for papers to creating the final program.



As the Pretalx API was completely changed, Pytanis needs an update.

Pyomo / HiGHS notebook:

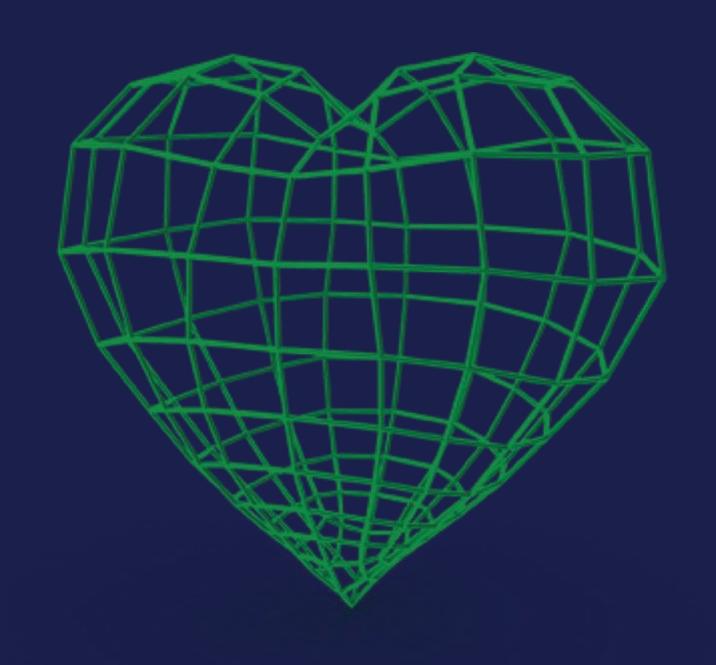
https://github.com/PioneersHub/pytanis/blob/main/notebooks/pyconde-pydata-darmstadt-2025/40_scheduling_v1.ipynb



Summary

References

- Introduction to Optimization by Laurent Lessard
- Mixed Integer Programming for time table scheduling by Francisco Espiga
- Schedule Optimisation using Linear Programming in Python by Lewis Woolfson
- Some icons taken from Flaticon.com





EuroScipy 2025

Thank you!







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